

Implementation of Randomized Dining Philosophers

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Outline



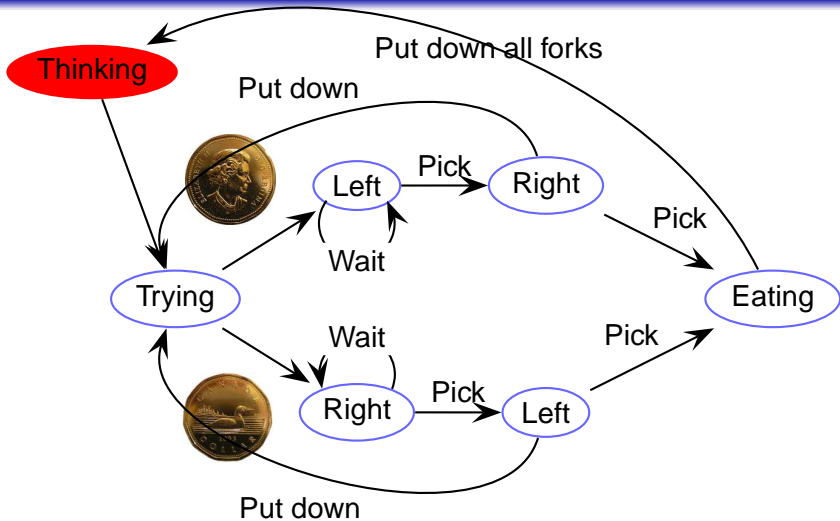
- 1 Introduction
- 2 Lehman and Rabin's Randomized Algorithm
- 3 Randomized Message Passing Algorithm
- 4 Conclusion

Recall

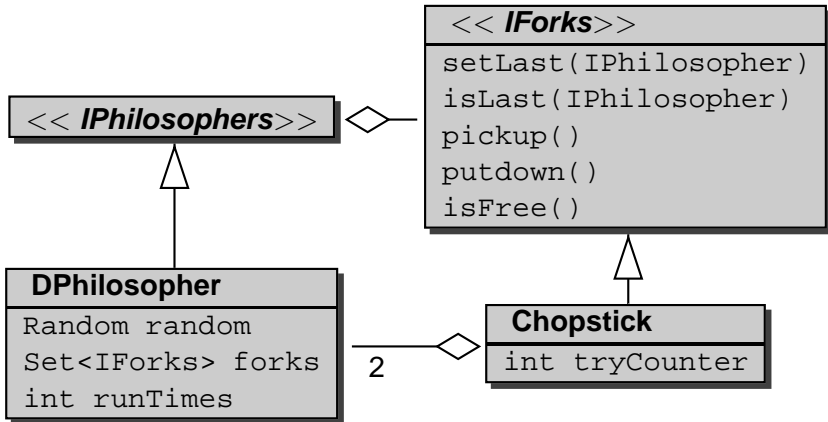


- We use randomized algorithm to break the initial symmetry
- Use different mechanism to guarantee mutual exclusion and deadlock free
- With extra information, we can achieve liveness property-starvation free

Control Flow Graph



Structure



Chopstick Object



```
1 public Semaphore getSemaphore() {
2     return semaphore;
3 }
4
5 public void pickup() {
6     state = INUSE;  continousTries = 0;
7 }
8 public boolean isLast(IPhilosopher last) {
9     if (last != this.last) return false;
10    else if (continousTries >= MAX_CONTINUOUS_TRY)
11        return false;
12    else {
13        continousTries ++;
14        return true;
15    }
```

DPhilosopher Object



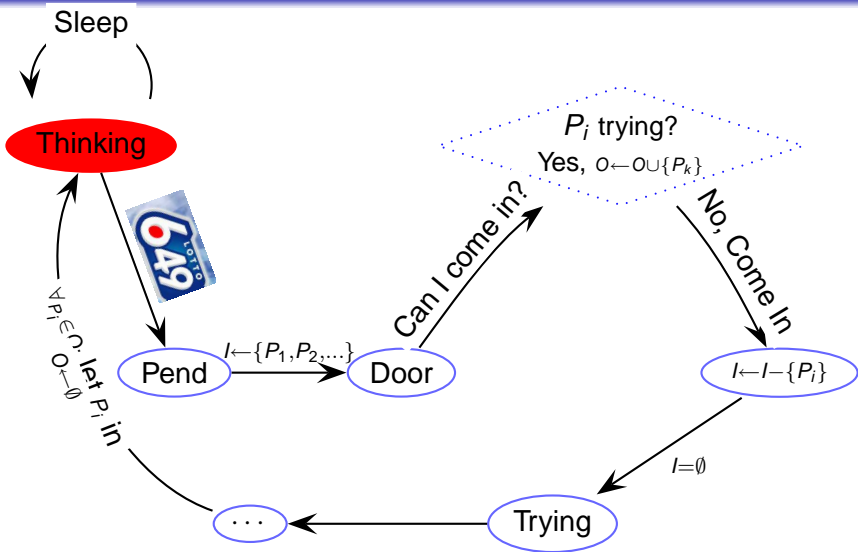
```
1  while(true){
2      if(random.nextInt(1) == 0){
3          try1 = left; try2 = right; }
4      else {
5          try1 = right; try2 = left; }
6
7      while{retry}{
8          try1.getSemaphore().acquire();
9          if (try1.isFree() && ! try1.isLast(this)){
10             try1.pickup(); retry = false;
11             try1.getSemaphore().release(); }
12         else{
13             try1.getSemaphore().release(); }
14     try2.getSemaphore().acquire();
15     if (try2.isFree()){
16         try2.pickup(); try2.getSemaphore().release
            (); }
```

DPhilosopher Object-cont

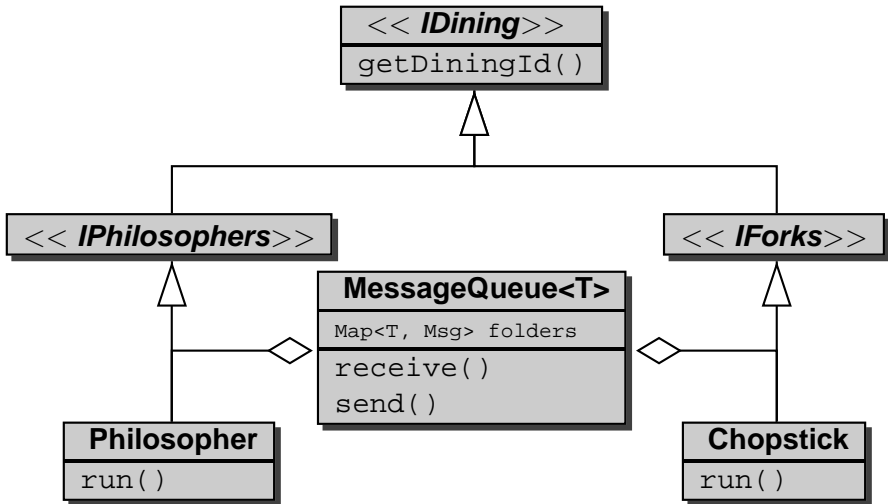


```
16
17     else { try2 .getSemaphore () .release () ;
18           try1 .getSemaphore () .acquire () ; try1 .putdown
19           () ;
20           try1 .getSemaphore () .release () ; continue ;
21     }
22     // eating
23     try1 .getSemaphore () .acquire () ;
24     try1 .putdown () ; try1 .setLast (this) ;
25     try1 .getSemaphore () .release () ;
26     try2 .getSemaphore () .acquire () ;
27     try2 .putdown () ; try2 .setLast (this) ;
28     try2 .getSemaphore () .release () ;
29 }
```


Control Flow



Structure



MessageQueue Object



```
1 public void send(IMessage<T> message) {
2     List<IMessage<T>> list = null;
3     T from = message.getRecipient();
4     synchronized(mFolders){
5         if ( !mFolders.containsKey(from) ){
6             list = Collections.synchronizedList(new
7                 LinkedList<IMessage<T>>());
8             mFolders.put(from, list); }
9         else {
10            list = mFolders.get(from); } }
11 }
12 public IMessage<T> receive(T recipient) {
13     synchronized(mFolders) {...}}
14 public IMessage<T> receiveBlocked(T recipient)
15     {...}
```

Chopstick Object



```
1 state = FREE;
2 while(true){
3     IMessage<IDining> msg = mQueue.receiveBlocked(
4         this);
5     if (type == TYPE_REGISTER){...}
6     else if (type == TYPE_REQUEST){...}
7     else if (type == TYPE_FAIL || type ==
8         TYPE_RELEASE){...}
9     else if (type == TYPE_REMOVE){...}
10 }
```

Philosopher Object



```
1  if (state == STATE_THINKING){
2    state = STATE_PENDING; in = new HashSet<IDining
      >();
3  for (IDining p : contenders){
4    if (p != this) { //send p request to enter door }
5  else if (state == STATE_TRYING && isTimeout()){
6    if (random.nextInt(contenders.size()) != 0){ //
      lottery
7    waitTime = (waitTime * 2) % MAX_WAIT_TIME }
8  else {
9    for(IForks fork : forks){//send pickup to
      forks}
10   state = STATE_HOPEFUL; grantedForkSet.clear();
      } }
11  else if (state == STATE_PENDING && in.isEmpty()){
12    state = STATE_TRYING;}
13 }
```

Philosopher Object



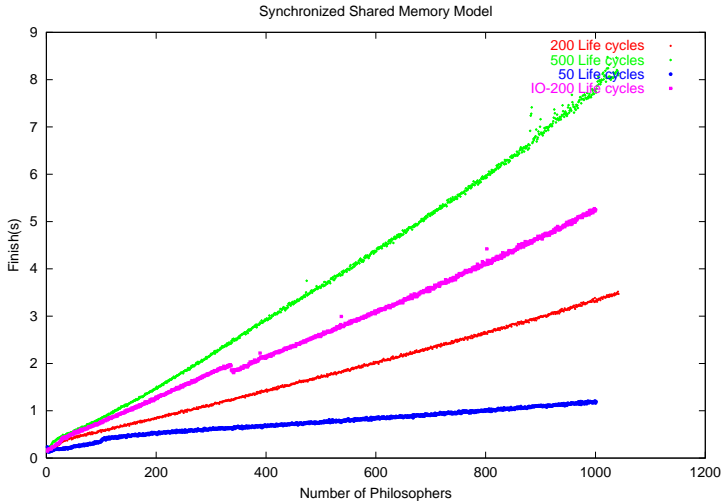
```
13 IMessage<IDining> msgIn = mQueue.receive(this);
14 IDining process = msgIn.getSender(); int type =
    msgIn.getType();
15 if (type == TYPE_UPDATE) {... }
16 else if (type == TYPE_ENTER){ //request from
    other philosophers to enter
17     if (state == STATE_TRYING ){ out.add(process); }
18     else{//send back to allow in} }
19 else if (type == TYPE_OUTSIDE && state ==
    STATE_PENDING){ //get permission to enter
20     in.remove(process); }
21 else if (type == TYPE_REMOVE) {... }
```

Philosopher Object

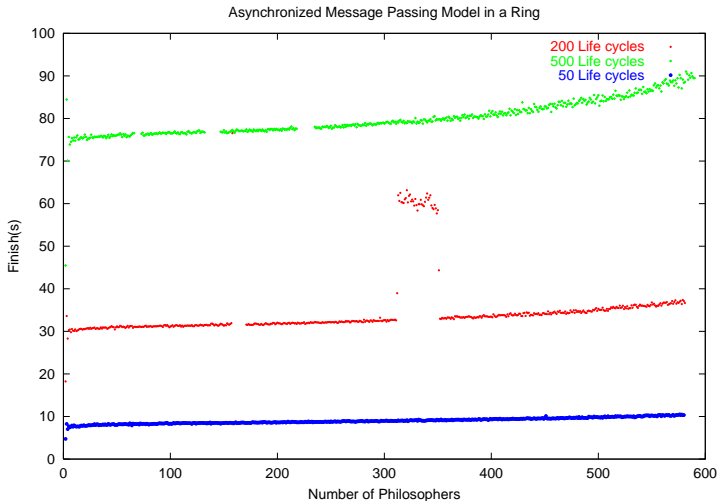


```
21 else if (type == TYPE_GRANT && state ==
    STATE_HOPEFUL){
22     grantedForkSet.add(process);
23     if (grantedForkSet.size() == forks.size()){
24         state = STATE_EATING;
25         //release fork by sending message
26         for (IDining outP: out){ //send outP allowance
            message } }
27 else if (type == TYPE_REJECT && state ==
    STATE_HOPEFUL){
28     for (IDining fork : grantedForkSet){//send fork
        failed message to release}
29     state = STATE_TRYING;
30     waitTime = (waitTime * 2) % MAX_WAIT_TIME;
31 else if (type == TYPE_GRANT && state !=
    STATE_HOPEFUL){//release }
```

Result for Shared Memory Algorithm



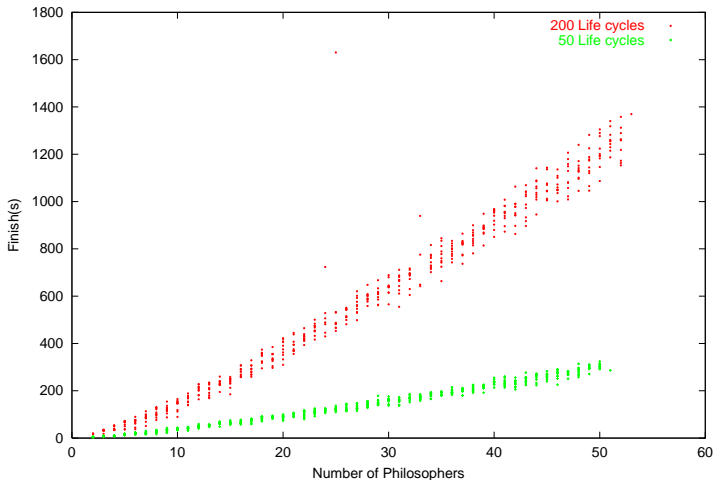
Result for Message Passing Algorithm in the Ring



Result for Message Passing Algorithm in the arbitrary graph



Asynchronized Message Passing Model in Arbitrary Graphs



Other Thoughts



- Both achieve linear or approximately linear time and space complexity regarding the number of thread
- Messageing passing algorithm is much slower and bigger varaince since it is asynchronized, and has more states
- Degree of contending affects the complexity of running time dramatically
- However messageing passing increases the flexibility
- It seems liveness property achieved, still need to verify

Future Work



- Expand synchronized shared memory model to arbitrary graphs
- Test different settings to find the real impact of graph and degree contending for asynchronized message passing model
- Try different mutual exclusion mechanism to increase performance in asynchronized message passing model