# Introduction to Functional Programming and basic Lisp 

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## Functional vs Declarative Programming

- declarative programming uses logical statements to describe objects.
* Prolog is an example of this kind of language.
- functional programming uses mathematical functions and functional expressions to describe objects.
* Functional programming has its roots in lambda calculus
* Lisp is a functional language.


## Functional Programming

- Functional programming is not based on assignments that change the state.
- Functions specify other values in terms of existing data without changing it.
- This allows all sorts of clever implementations e.g. on parallel hardware.


## Background on LISP

- Acronym for LISt Processing
- created by the AI pioneer John McCarthy
- widely used in research on AI for over 40 years.

ภUsed in industry to develop expert systems \& other AI applications

- found inside applications like Emacs and AutoCAD as an embedded language

ภmakes the embedding application easily extensible

## Lisp as an extensional language

- embedding Lisp makes it easy to extend an application
- creation of new languages built "on" Lisp
- industrial-strength versions are usually standardized on Common Lisp


## LISP Interpreter

- An interactive environment, always evaluating input
- To run LISP at prism labs, type:
"clisp"
To exit: Ctrl+D
- To load the file Ip.Isp, in C: \MyFolder, use the following command: (load "C://MyFolder/lp.lsp")


## S-expressions

- A symbolic expression (s-expression) is defined inductively as
* an atom (number or word), or
* dotted pair of s-expressions (e.g. (x.y) where $x$ and $y$ are s-expressions)
- lists are the main kind of s-expressions
* Example: (a b c) or a. (b . (c. nil)) - nil is the same as () - the only atom that is a list


## disassembly

- functions which extract the two parts of a dotted pair:
* first extracts the first part,凤also called car
* rest extracts the second part, дalso called cdr


## contents of a dotted-pair



## new names for old

- car = function which returns the content referenced by address-register
- cdr = function which returns content referenced by decrement register
- many books cling to car and cdr for 'backwards compatibility'.
* can use first and rest
* you can use either - the Lisp interpreter doesn't care.


## list structure



## functional expressions

- terms (functional expressions) are represented as lists
$\delta_{\text {w }}$ write $f(x, y)$ as ( $f x y$ ).
$\Omega(a b c)$ represents the term $a(b, c)$.
- already we see a bit of the power of symbolic computing:
* expressions have same form as data
* a function name is just an atom (a symbol)
* could itself be computed


## evaluating functions

- evaluating (function arg1 arg2 . . )
* applies the system function eval to each argument
ภthen applies the function to the results
* $(+2(+35))$
eval(2) = 2
$\operatorname{eval}((+35))=+($ eval $(3), \operatorname{eval}(5))=+(3,5)=8$
$+(2,8)=10$


## blocking evaluation

$\bullet$ try evaluating

* (reverse (a b)) --> error $\Omega$ no function a defined
* (reverse (+ 12 )) --> error
$\Omega^{\prime \prime} 3$ " is not a list
- how to block evaluation of an sexpression?
* quote it


## the uses of quotation

- try evaluating
* (reverse (quote (a b))) --> (B A)

ภquote returns its argument unevaluated

* (reverse '(+ 12$)$ ) --> (2 1 +)
$\Omega^{\prime}<s$-expression> $=$ (quote <s-expression>)
- (equal '(reverse (a b)) '(b a))) ?
* NIL -- why?
- (equal (reverse '(reverse (a b))) '((a b) reverse)))
* T


## List Processing Functions

$$
\begin{aligned}
& >\left(\operatorname{car}^{\prime}(\text { a b c })\right)-->\text { A } \\
& >\left(\text { cdr }^{\prime}(\text { a b c) })-->(\text { B C) }\right. \\
& >\left(\operatorname{car}\left(\operatorname{cdr}\left(\operatorname{car}^{\prime}((\operatorname{ab}))\right)\right)\right)-->\text { (B) }
\end{aligned}
$$

* Other built-in functions for list processing include: cons, append, list, ...
$>($ cons 'a '(b c)) $-->(A B C)$
$>$ (append $\left.{ }^{\prime}(\mathrm{a})^{\prime}(\mathrm{b})^{\prime}(\mathrm{c})\right)-->(\mathrm{A} \mathrm{B} \mathrm{C})$
> (list `a 'b 'c) --> (A B C)


## oceans of functions

- the basis for Lisp programs is the concept of a * and variants with special properties
- Common Lisp has several hundred built-in functions
- many are redundant - could be replaced by expressions involving other functions


## Function definitions

- Written as (defun function-name arg-list result-spec).
- function-name is a symbol.
- arg-list is a list of symbols, the parameters.
- result-spec is an expression whose value is the result of the function.
- When a function application is evaluated, substitute actual arguments for parameters in result-spec and return its value.


## Function Example 1

$>($ defun avg (xy) (/ (+ x y) 2.0))
AVG
$>(\operatorname{avg} 12)$
1.5

## Function Example 2

- Rewrite (John P Doe) as (Doe John P)
- (defun last_name_first (name_list) (cons (third name_list)
(cons (first name_list)
(cons (second name_list) nil))))


## Functions in LISP

- functions are considered as first class objects in Lisp.
- A function can take another function as an argument and a function can return a function as a value.
- This is what make functional programming very powerful.
- In a sense you can define your own control structures and manipulate programs!


## More

- 'Practical Common Lisp' Book online http://www.gigamonkeys.com/book/
- Notes on Lambda Calculus http://www.mathstat.dal.ca/~selinger/papers/la mbdanotes.pdf

