# Path Testing and Test Coverage

Chapter 9

## Structural Testing

- Also known as glass/white/open box testing
- Structural testing is based on using specific knowledge of the program source text to define test cases
  - Contrast with functional testing where the program text is not seen but only hypothesized

## Structural Testing

- Structural testing methods are amenable to
  - Rigorous definitions
    - Control flow, data flow, coverage criteria
  - Mathematical analysis
    - Graphs, path analysis
  - Precise measurement
    - Metrics, coverage analysis



## • What is a program graph?

## Program Graph Definition – 2

- Given a program written in an imperative programming language
  - Its program graph is a directed graph in which nodes are statements and statement fragments, and edges represent flow of control
  - Two nodes are connected if execution can proceed from one to the other



```
1 output ("Enter 3 integers")
```

```
2 input (a, b, c)
```

```
3 output("Side a b c: ", a, b, c)
```

```
4 if (a < b) and (b < a+c) and (c < a+b)
```

```
5 then isTriangle ← true
```

```
6 else isTriangle ← false
```

fi

```
7 fi
```

```
8 if isTriangle
```

```
9 then if (a = b) and (b = c)
```

```
10 else output ("equilateral")
```

```
else if (a \neq b) and (a \neq c) and (b \neq c)
```

```
12 then output ("scalene")
```

```
13 else output("isosceles")
```

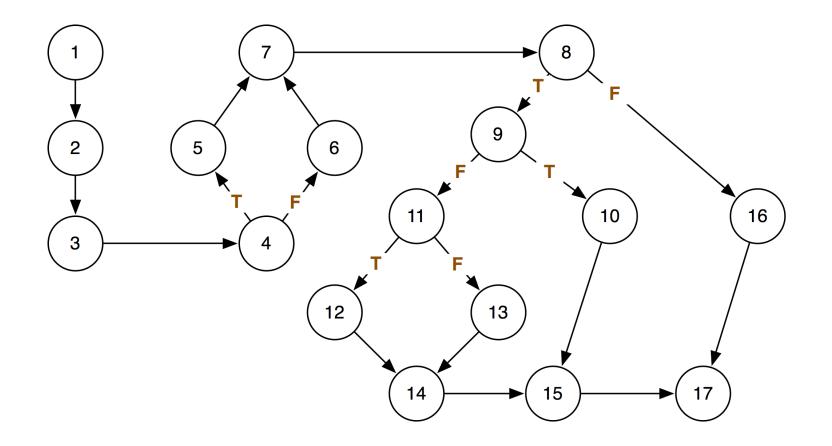
```
14
```

```
15 fi
```

```
16 else output ("not a triangle")
```

```
17 fi
```



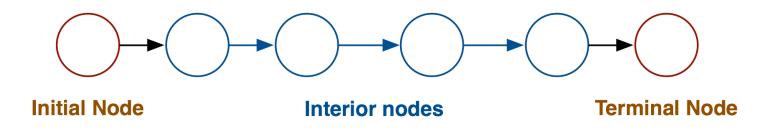




## • What is a DD-path?

## DD-Path – informal definition

- A decision-to-decision path (DD-Path) is a path chain in a program graph such that
  - Initial and terminal nodes are distinct
  - Every interior node has indeg = 1 and outdeg = 1
    - The initial node is 2-connected to every other node in the path
    - No instances of 1- or 3-connected nodes occur





## What is the definition of node connectedness?

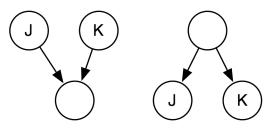
Hint: There are 4-types of connectedness



- Two nodes J and K in a directed graph are
  - 0-connected iff no path exists between them

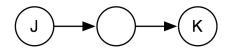


 1-connected iff a semi-path but no path exists between them

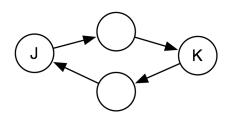


## Connectedness definition – 2

- Two nodes J and K in a directed graph are
  - 2-connected iff a path exists between between them



- 3-connected iff a path goes from J to K , and a path goes from K to  $n_1$ 



## DD-Path – formal definition

- A decision-to-decision path (DD-Path) is a chain in a program graph such that:
  - Case 1: consists of a single node with indeg=0
  - Case 2: consists of a single node with outdeg=0
  - Case 3: consists of a single node with indeg ≥ 2 or outdeg ≥ 2
  - Case 4: consists of a single node with indeg =1, and outdeg = 1
  - Case 5: it is a maximal chain of length  $\geq 1$
- DD-Paths are also known as segments

# Triangle program DD-paths

Nodes	Path	Case
1	First	1
2,3	А	5
4	В	3
5	С	4
6	D	4
7	E	3
8	F	3
9	G	3

Nodes	Path	Case
10	Н	4
11	Ι	3
12	J	4
13	К	4
14	L	3
15	М	3
16	N	4
17	Last	2



## • What is a DD-path graph?

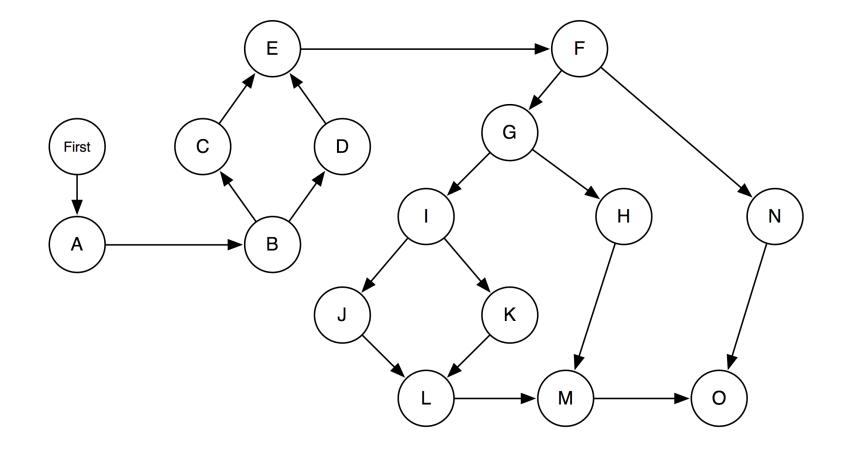
## DD-Path Graph – informal definition

- Given a program written in an imperative language, its
   DD-Path graph is a directed graph, in which
  - Nodes are DD-Paths of its program graph
  - Edges represent control flow between successor DD-Paths.
- Also known as Control Flow Graph

## **Control Flow Graph Derivation**

- Straightforward process
- Some judgment is required
- The last statement in a segment must be
  - a predicate
  - a loop control
  - a break
  - a method exit







```
public int displayLastMsg(int nToPrint) {
  np = 0;
  if ((msgCounter > 0) && (nToPrint > 0)) {
    for (int j = lastMsg; (( j != 0) && (np < nToPrint)); --j) {</pre>
      System.out.println(messageBuffer[j]);
      ++np;
    }
    if (np < nToPrint) {</pre>
      for (int j = SIZE; ((j != 0) && (np < nToPrint)); --j) {</pre>
        System.out.println(messageBuffer[j]);
        ++np;
      }
  return np;
```

# displayLastMsg– Segments part 1

#### Line

#### Segment

1	<pre>public int displayLastMsg(int nToPrint) {</pre>	
2	np = 0;	Α
3a	if ( (msgCounter > 0)	A
3b	&& (nToPrint > 0))	В
4a	{ for (int j = lastMsg;	С
4b	( (j!= 0)	D
4c	&& (np < nToPrint));	Е
4d	j)	F
5	<pre>{ System.out.println(messageBuffer[j]);</pre>	F
6	++np;	F
7	}	F

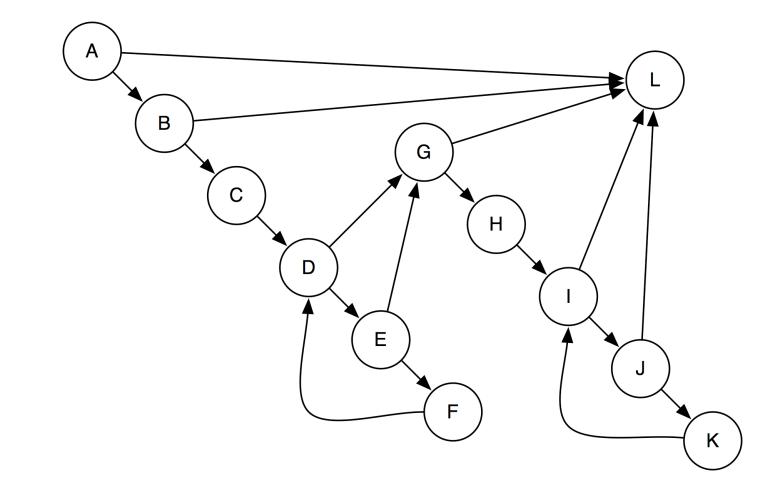
# displayLastMsg– Segments part 2

#### Line

#### Segment

8	if (np < nToPrint)	G
9a	{ for (int j = SIZE;	Н
9b	((j != 0) &&	Ι
9c	(np < nToPrint));	J
9d	j)	К
10	<pre>{ System.out.println(messageBuffer[j]);</pre>	К
11	++np;	К
12	}	L
13	}	L
14	}	L
15	return np;	L
16	}	L

displayLastMsg – Control Flow Graph



PT-22

## Control flow graphs definition – 1

- Depict which program segments may be followed by others
- A segment is a node in the CFG
- A conditional transfer of control is a branch represented by an edge
- An entry node (no inbound edges) represents the entry point to a method
- An exit node (no outbound edges) represents an exit point of a method

## Control flow graphs definition – 2

- An entry-exit path is a path from the entry node to the exit node
- Path expressions represent paths as sequences of nodes
- Loops are represented as segments within parentheses followed by an asterisk
- There are 22 different entry-exit path expressions in displayLastMsg

## Entry-exit path expressions – part 1

#### Entry-Exit paths

1	AL
2	ABL
3	ABCDGL
4	ABCDEGL
5	A B C (D E F)* D G L
6	A B C (D E F)* D E G L
7	ABCDGHIL
8	ABCDGHIJL
9	A B C D G H (I J K)* I L
10	A B C (D E F)* D E G H (I J K)* I J L
11	ABCDEGHIL

Entry-exit path expressions – part 2

#### Entry-Exit paths

12	ABCDEGHIJL
13	A B C D E G H (I J K)* I L
14	A B C D E G H (I J K)* I J L
15	ABC (DEF)* DGHIL
16	A B C (D E F)* D G H I J L
17	A B C (D E F)* D G H (I J K)* I L
18	A B C (D E F)* D G H (I J K)* I J L
19	ABC (DEF)* DEGHIL
20	ABC (DEF)* DEGHIJL
21	A
22	A B C (D E F)* D E G H (I J K)* I J L

## Paths displayLastMsg – decision table – part 1

#### Path condition by Segment Name

	Entry/Exit Path	А	В	D	Е	G	Ι	J
1	AL	F	-	-	_	—	_	—
2	ABL	Т	F	-	_	—	_	—
3	ABCDGL	Т	Т	F	_	F	-	_
4	ABCDEGL	Т	Т	Т	F	_	-	_
5	A B C (D E F)* D G L	Т	Т	T/F	T/-	F	-	_
6	A B C (D E F)* D E G L	Т	Т	T/T	T/F	F	-	_
7	ABCDGHIL	Т	Т	F	_	Т	F	_
8	ABCDGHIJL	Т	Т	F	_	Т	Т	F
9	A B C D G H (I J K)* I L	Т	Т	F	_	T/F	T/-	Т
10	A  B  C  D  G  H  (I  J  K)* I  J  L	Т	Т	F	_	T/T	T/F	Т
11	ABCDEGHIL	Т	Т	Т	F	Т	F	_

x/x Conditions at loop-entry / loop-exit – is don't care

PT-27

## Paths displayLastMsg – decision table – part 2

#### Path condition by Segment Name

	Entry/Exit Path	А	В	D	Е	G	Ι	J
12	ABCDEGHIJL	Т	Т	Т	F	Т	Т	F
13	A B C D E G H (I J K)* I L	Т	Т	Т	F	Т	T/F	T/-
14	A	Т	Т	Т	F	Т	T/T	T/F
15	A B C (D E F)* D G H I L	Т	Т	T/F	T/-	Т	F	_
16	A B C (D E F)* D G H I J L	Т	Т	T/T	T/F	Т	Т	F
17	A B C (D E F)* D G H (I J K)* I L	Т	Т	T/F	T/-	Т	T/F	T/-
18	A B C (D E F)* D G H (I J K)* I J L	Т	Т	T/F	T/-	Т	T/T	T/F
19	A B C (D E F)* D E G H I L	Т	Т	T/T	T/F	Т	F	_
20	A B C (D E F)* D E G H I J L	Т	Т	T/T	T/F	Т	Т	F
21	A B C (D E F)* D E G H (I J K)* I L	Т	Т	T/T	T/F	Т	Т	Т
22	A B C (D E F)* D E G H (I J K)* I J L	Т	Т	T/T	T/F	Т	Т	Т

x/x Conditions at loop-entry / loop-exit – is don't care



## • List the program text coverage metrics.

## Program text coverage Metrics – 2

- C<sub>0</sub> Every Statement
- C<sub>1</sub> Every DD-path
- C<sub>1p</sub> Every predicate to each outcome
- $C_2$   $C_1$  coverage + loop coverage
- $C_d$   $C_1$  coverage + every dependent pair of DD-paths
- C<sub>MCC</sub> Multiple condition coverage
- C<sub>ik</sub> Every program path that contains k loop repetitions
- C<sub>stat</sub> Statistically significant fraction of the paths
- $C_{\infty}$  Every executable path



# What are the common program text coverage models?



- Statement Coverage
- Segment Coverage
- Branch Coverage
- Multiple-Condition Coverage



## • When is statement coverage achieved?

## Statement coverage $-C_0 - 2$

- Achieved when all statements in a method have been executed at least once
- A test case that will follow the path expression below will achieve statement coverage in our example

# **A B C (D E F)\* D G H (I J K)\* I L**

• One test case is enough to achieve statement coverage!



• When is segment coverage achieved?

## Segment coverage – 2

- Achieved when all segments have been executed at least once
  - Segment coverage counts segments rather than statements
  - May produce drastically different numbers
    - Assume two segments P and Q
    - P has one statement, Q has nine
    - Exercising only one of the segments will give either 10% or 90% statement coverage
    - Segment coverage will be 50% in both cases



### • What problems are there with statement coverage?

Statement coverage problems – 2

- Important cases may be missed
  - Predicate may be tested for only one value
    - misses many bugs
  - Loop bodies may only be iterated only once

String s = null; if (x != y) s = "Hi"; String s2 = s.substring(1);

- What coverage solves this problem?
  - Define it



- Achieved when every edge from a node is executed at least once
- At least one true and one false evaluation for each predicate

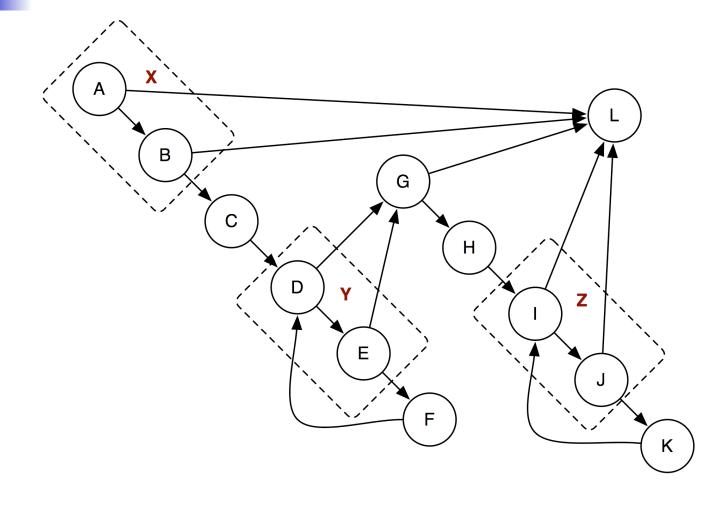
How many test cases are required?

# Branch coverage $-C_{1p} - 2$

- Can be achieved with D+1 paths in a control flow graph with D 2-way branching nodes and no loops
  - Even less if there are loops
- In the Java example displayLastMsg branch coverage is achieved with three paths – see next few slides

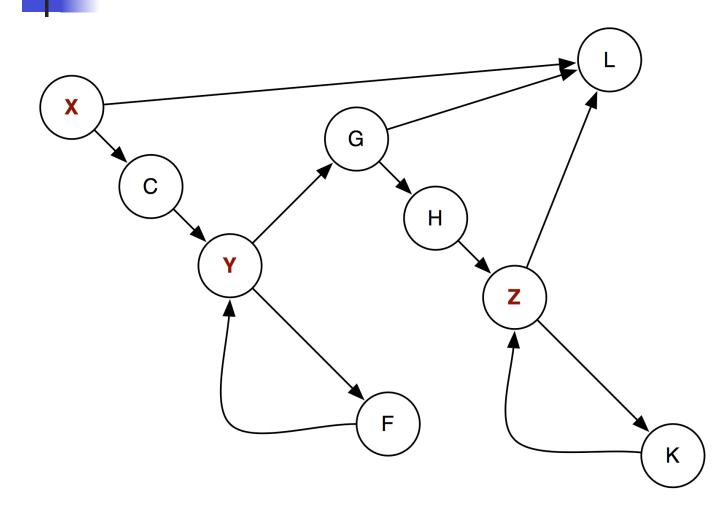
```
X L
X C (Y F)* Y G L
X C (Y F)* Y G H (Z K)* Z L
```

### Java example program displayLastMsg – DD-path graph



X, Y & Z are shorthand for the nodes within the dotted boxes; used for branch testing

## Java example program displastLastMsg – aggregate predicate DD-path graph



# Aggregate Paths – decision table – part 1

#### Path condition by Segment Name

	Branch Coverage	А	В	D	Е	G	Ι	J
1	XL	F	_	_	_	_	_	_
2	XL	Т	F	-	_	—	_	—
3	XCYGL	Т	Т	F	_	F	-	—
4	XCYGL	Т	Т	Т	F	_	-	—
5	X C (Y F)* Y G L	Т	Т	T/F	T/-	F	-	—
6	X C (Y F)* Y G L	Т	Т	T/T	T/F	F	-	—
7	XCYGHZL	Т	Т	F	_	Т	F	—
8	XCYGHZL	Т	Т	F	_	Т	Т	F
9	X C Y G H (Z K)* I L	Т	Т	F	_	T/F	T/-	Т
10	X C Y G H (Z K)* I L	Т	Т	F	-	T/T	T/F	Т
11	XCYGHZL	Т	Т	Т	F	Т	F	—

x/x Conditions at loop-entry / loop-exit – is don't care

# Aggregate Paths – decision table example – part 2

	Branch Coverage	А	В	D	Е	G	Ι	J
12	XCYGHZL	Т	Т	Т	F	Т	Т	F
13	X C Y G H (Z K)* Z L	Т	Т	Т	F	Т	T/F	T/-
14	X C Y G H (Z K)* Z L	Т	Т	Т	F	Т	T/T	T/F
15	X C (Y F)* Y G H Z L	Т	Т	T/F	T/-	Т	F	_
16	X C (Y F)*Y G H Z L	Т	Т	T/T	T/F	Т	Т	F
17	X C (Y F)* Y G H (Z K)* Z L	Т	Т	T/F	T/-	Т	T/F	T/-
18	X C (Y F)* Y G H (Z K)* Z L	Т	Т	T/F	T/-	Т	T/T	T/F
19	X C (Y F)* Y G H Z L	Т	Т	T/T	T/F	Т	F	_
20	X C (Y F)* Y G H Z L	Т	Т	T/T	T/F	Т	Т	F
21	X C (Y F)* Y G H (Z K)* Z L	Т	Т	T/T	T/F	Т	Т	Т
22	X C (Y F)* Y G H (Z K)* Z L	Т	Т	T/T	T/F	Т	Т	Т

#### Path condition by Segment Name

x/x Conditions at loop-entry / loop-exit — is don't care



### • What are the problems with branch coverage?



- Ignores implicit paths from compound paths
  - 11 paths in aggregate model vs 22 in full model

### Branch coverage problems – 3

- Ignores implicit paths from compound paths
  - 11 paths in aggregate model vs 22 in full model
- Short-circuit evaluation means that many predicates might not be evaluated
  - A compound predicate is treated as a single statement. If n clauses, 2<sup>n</sup> combinations, but only 2 are tested

### Branch coverage problems – 4

- Ignores implicit paths from compound paths
  - 11 paths in aggregate model vs 22 in full model
- Short-circuit evaluation means that many predicates might not be evaluated
  - A compound predicate is treated as a single statement. If n clauses, 2<sup>n</sup> combinations, but only 2 are tested

### Only a subset of all entry-exit paths is tested

Two tests for branch coverage vs 4 tests for path coverage

• 
$$a = b = x = y = 0$$
 and  $a = x = 0 \land b = y = 1$ 

Overcoming branch coverage problems

How do we overcome branch coverage problems?

# Overcoming branch coverage problems – 2

### Use Multiple condition coverage

- All true-false combinations of simple conditions in compound predicates are considered at least once
  - Guarantees statement, branch and predicate coverage
  - Does not guarantee path coverage
- A truth table may be necessary
- Not necessarily achievable
  - lazy evaluation true-true and true-false are impossible
  - mutually exclusive conditions false-false branch is impossible

if ((x > 0) || (x < 5)) ...

## Overcoming branch coverage problems – 3

- Can have infeasible paths due to dependencies and redundant predicates
  - Paths perpetual .. motion and free .. lunch are impossible
  - In this case indicates a potential bug
    - At least poor program text

```
if x = 0 then oof.perpetual
        else off.free
fi

if x != 0 then oof.motion
        else off.lunch
fi
```



- Loops are highly fault-prone, so they need to be tested carefully
- Based on the previous slides on testing decisions what would be a simple view of testing a loop?

### Dealing with Loops – 2

- Simple view
  - Involves a decision to traverse the loop or not
  - Test as a two way branch
- What would functional testing suggest as a better way of testing?
- What tests does it suggest?

### Dealing with Loops – 3

- A bit better
  - Boundary value analysis on the index variable
  - Suggests a zero, one, many tests

### How do we deal with nested loops?



- Nested loops
  - Tested separately starting with the innermost
- Once loops have been tested what can we do with the control flow graph?



- Once loops have been tested
  - They can be condensed to a single node

### Condensation graphs

- Condensation graphs are based on removing strong components or DD-paths
- For programs remove structured program constructs
  - One entry, one exit constructs for sequences, choices and loops
  - Each structured component once tested can be replaced by a single node when condensing its graph

### Violations of proper structure

- Program text that violates proper structure cannot be condensed
  - Branches either into or out of the middle of a loop
  - Branches either into or out of then and else phrases of if... then...else statements
  - Increases the complexity of the program
  - Increases the difficulty of testing the program

### Cyclomatic number

- The cyclomatic number for a graph is given by
  - CN(G) = e v + 2\*c
    - e number of edges
       v number of vertices
       c number of connected regions
  - For strongly connected graphs, need to add edges from every sink to every source

## Cyclomatic number for programs

- For properly structured programs there is only one component with one entry and one exit. There is no edge from exit to entry.
- Definition 1: CN(G) = e v + 2
  - Only 1 component, not strongly connected
- Definition 2: CN(G) = p + 1
  - p is the number of predicate nodes with out degree = 2
- Definition 3: CN(G) = r + 1
  - r is the number of enclosed regions