11-711 Algorithms for NLP

The Earley Parsing Algorithm

Reading:

Jay Earley,

“An Efficient Context-Free Parsing Algorithm”

Comm. of the ACM vol. 13 (2), pp. 94–102

The Earley Parsing Algorithm

General Principles:

- A clever hybrid *Bottom-Up* and *Top-Down* approach
- *Bottom-Up* parsing completely guided by *Top-Down* predictions
- Maintains sets of “dotted” grammar rules that:
  - Reflect what the parser has “seen” so far
  - Explicitly predict the rules and constituents that will combine into a complete parse
- Similar to Chart Parsing - partial analyses can be shared
- Time Complexity $O(n^3)$, but better on particular sub-classes
- Developed prior to Chart Parsing, first efficient parsing algorithm for general context-free grammars.

The Earley Parsing Method

- Main Data Structure: The *"state"* (or *"item"*)
- A state is a "dotted" rule and starting position:
  $[A \rightarrow X_1 ... \bullet C ... X_m, p_i]$
- The algorithm maintains sets of "states", one set for each position in the input string (starting from 0)
- We denote the set for position i by S_i

The Earley Parsing Algorithm

Three Main Operations:

- **Predictor:** If state $[A \to X_1 ... \bullet C ... X_m, j] \in S_i$ then for every rule of the form $C \to Y_1 ... Y_k$, add to S_i the state $[C \to \bullet Y_1 ... Y_k, i]$
- **Completer:** If state $[A \to X_1 ... X_m \bullet, j] \in S_i$ then for every state in S_j of form $[B \to X_1 ... \bullet A ... X_k, l]$, add to S_i the state $[B \to X_1 ... A \bullet ... X_k, l]$
- **Scanner:** If state $[A \to X_1 ... \bullet a ... X_m, j] \in S_i$ and the next input word is $x_{i+1} = a$, then add to S_{i+1} the state $[A \to X_1 ... a \bullet ... X_m, j]$

The Earley Recognition Algorithm

- Simplified version with no lookaheads and for grammars without epsilon-rules
- Assumes input is string of grammar terminal symbols
- We extend the grammar with a new rule $S' \rightarrow S$ \$
- The algorithm sequentially constructs the sets S_i for $0 \leq i \leq n+1$
- We initialize the set S_0 with $S_0 = \{[S' \rightarrow \bullet S \ \$, 0]\}$

The Earley Recognition Algorithm

The Main Algorithm: parsing input $x = x_1 ... x_n$

1. $S_0 = \{[S' \rightarrow \bullet S \ \$, 0]\}$
2. For $0 \leq i \leq n$ do:
   Process each item $s \in S_i$ in order by applying to it the *single* applicable operation among:
   (a) Predictor (adds new items to S_i)
   (b) Completer (adds new items to S_i)
   (c) Scanner (adds new items to S_{i+1})
3. If $S_{i+1} = \phi$, *Reject* the input
4. If $i = n$ and $S_{n+1} = \{[S' \rightarrow S \ \$\bullet, 0]\}$ then *Accept* the input

Earley Recognition - Example

The Grammar:

$$
\begin{array}{rlcl}
(1) & S & \rightarrow & NP\,VP \\
(2) & NP & \rightarrow & art\ adj\ n \\
(3) & NP & \rightarrow & art\ n \\
(4) & NP & \rightarrow & adj\ n \\
(5) & VP & \rightarrow & aux\ VP \\
(6) & VP & \rightarrow & v\ NP
\end{array}
$$

The original input: " $x =$ The large can can hold the water"
POS assigned input: " $x =$ art adj n aux v art n"
Parser input: " $x =$ art adj n aux v art n \$"

Earley Recognition - Example

The input: "$x =$ **art** adj n aux v art n \$"

$$
\begin{aligned}
S_0\text{:}\quad & [S' \rightarrow \bullet S \text{ \$ },\ 0] \\
& [S \rightarrow \bullet NP\ VP\ ,\ 0] \\
& [NP \rightarrow \bullet art\ adj\ n\ ,\ 0] \\
& [NP \rightarrow \bullet art\ n\ ,\ 0] \\
& [NP \rightarrow \bullet adj\ n\ ,\ 0]
\end{aligned}
$$

$$
\begin{aligned}
S_1\text{:}\quad & [NP \rightarrow art\ \bullet adj\ n\ ,\ 0] \\
& [NP \rightarrow art\ \bullet n\ ,\ 0]
\end{aligned}
$$

Earley Recognition - Example

The input: "$x =$ art **adj** n aux v art n $"

S_1: $[NP \rightarrow art \ \bullet adj \ n \ , \ 0]$
$\quad\quad [NP \rightarrow art \ \bullet n \ , \ 0]$

S_2: $[NP \rightarrow art \ adj \ \bullet n \ , \ 0]$

Earley Recognition - Example

The input: "$x =$ art adj **n** aux v art n \$"

S_2: $[NP \rightarrow art\ adj\ \bullet n\ ,\ 0]$

S_3: $[NP \rightarrow art\ adj\ n \bullet\ ,\ 0]$

Earley Recognition - Example

The input: “$x =$ art adj n **aux** v art n \$”

S_3: $[NP \rightarrow art\ adj\ n \bullet,\ 0]$
$[S \rightarrow NP\ \bullet VP,\ 0]$
$[VP \rightarrow \bullet aux\ VP,\ 3]$
$[VP \rightarrow \bullet v\ NP,\ 3]$

S_4: $[VP \rightarrow aux\ \bullet VP,\ 3]$

Earley Recognition - Example

The input: "$x =$ art adj n aux **v** art n \$"

S_4: $[VP \rightarrow aux\ \bullet VP\ ,\ 3]$

$[VP \rightarrow \bullet aux\ VP\ ,\ 4]$

$[VP \rightarrow \bullet v\ NP\ ,\ 4]$

S_5: $[VP \rightarrow v\ \bullet NP\ ,\ 4]$

Earley Recognition - Example

The input: "$x =$ art adj n aux v **art** n \$"

S_5: $[VP \to v \bullet NP , 4]$
$[NP \to \bullet art\ adj\ n , 5]$
$[NP \to \bullet art\ n , 5]$
$[NP \to \bullet adj\ n , 5]$

S_6: $[NP \to art \bullet adj\ n , 5]$
$[NP \to art \bullet n , 5]$

Earley Recognition - Example

The input: "$x =$ art adj n aux v art **n** \$"

S_6: $[NP \rightarrow art\ \bullet adj\ n\ ,\ 5]$
$[NP \rightarrow art\ \bullet n\ ,\ 5]$

S_7: $[NP \rightarrow art\ n \bullet\ ,\ 5]$

Earley Recognition - Example

The input: "$x =$ art adj n aux v art n **$**"

S_7: $[NP \rightarrow art\ n\ \bullet\ ,\ 5]$
$[VP \rightarrow v\ NP\ \bullet\ ,\ 4]$
$[VP \rightarrow aux\ VP\ \bullet\ ,\ 3]$
$[S \rightarrow NP\ VP\ \bullet\ ,\ 0]$
$[S' \rightarrow S\ \bullet\ \$\ ,\ 0]$

S_8: $[S' \rightarrow S\ \$\ \bullet\ ,\ 0]$

Time Complexity of Earley Algorithm

- Algorithm iterates for each word of input (i.e. n iterations)
- How many items can be created and processed in S_i?
  - Each item in S_i has the form $[A \rightarrow X_1 ... \bullet C ... X_m, j]$, $0 \leq j \leq i$
  - Thus $O(n)$ items
- The *Scanner* and *Predictor* operations on an item each require constant time
- The *Completer* operation on an item adds items of form $[B \rightarrow X_1 ... A \bullet ... X_k, l]$ to S_i, with $0 \leq l \leq i$, so it may require up to $O(n)$ time for each processed item
- Time required for each iteration (S_i) is thus $O(n^2)$
- Time bound on entire algorithm is therefore $O(n^3)$

Time Complexity of Earley Algorithm

Special Cases:

- *Completer* is the operation that may require $O(i^2)$ time in iteration i
- For unambiguous grammars, Earley shows that the completer operation will require at most $O(i)$ time
- Thus time complexity for unambiguous grammars is $O(n^2)$
- For some grammars, the number of items in each S_i is bounded by a *constant*
- These are called *bounded-state* grammars and include even some ambiguious grammars.
- For bounded-state grammars, the time complexity of the algorithm is linear - $O(n)$

Parsing with an Earley Parser

- As usual, we need to keep back-pointers to the constituents that we combine together when we complete a rule
- Each item must be extended to have the form $[A \rightarrow X_1(pt_1)... \bullet C...X_m, j]$, where the pt_i are “pointers” to the already found RHS sub-constituents
- At the end - reconstruct parse from the “back-pointers”
- To maintain efficiency - we must do ambiguity packing

Earley Parsing - Example

The input: “$x =$ art adj n aux v art n \$”

Earley Parsing - Example

The input: "$x =$ **art** adj n aux v art n \$"

S_0: $[S' \rightarrow \bullet S \,\$\,,\ 0]$
$[S \rightarrow \bullet NP\ VP\ ,\ 0]$
$[NP \rightarrow \bullet art\ adj\ n\ ,\ 0]$
$[NP \rightarrow \bullet art\ n\ ,\ 0]$
$[NP \rightarrow \bullet adj\ n\ ,\ 0]$

S_1: $[NP \rightarrow art_1\ \bullet adj\ n\ ,\ 0]$ 1 art
$[NP \rightarrow art_1\ \bullet n\ ,\ 0]$

Earley Parsing - Example

The input: "$x =$ art **adj** n aux v art n \$"

S_1: $[NP \rightarrow art_1 \bullet adj\ n\ ,\ 0]$
$\quad [NP \rightarrow art_1 \bullet n\ ,\ 0]$

S_2: $[NP \rightarrow art_1\ adj_2 \bullet n\ ,\ 0]$ 2 adj

Earley Parsing - Example

The input: "$x =$ art adj **n** aux v art n \$"

S_2: $[NP \rightarrow art_1\ adj_2\ \bullet n\ ,\ 0]$

S_3: $[NP_4 \rightarrow art_1\ adj_2\ n_3 \bullet\ ,\ 0]$

3 n

4 $NP \rightarrow art_1\ adj_2\ n_3$

Earley Parsing - Example

The input: "$x =$ art adj n **aux** v art n \$"

S_3: $[NP_4 \rightarrow art_1\ adj_2\ n_3 \bullet,\ 0]$
$[S \rightarrow NP_4\ \bullet VP,\ 0]$
$[VP \rightarrow \bullet aux\ VP,\ 3]$
$[VP \rightarrow \bullet v\ NP,\ 3]$

S_4: $[VP \rightarrow aux_5\ \bullet VP,\ 3]$ 5 aux

Earley Parsing - Example

The input: "$x =$ art adj n aux **v** art n \$"

S_4: $[VP \rightarrow aux_5 \ \bullet VP \ , \ 3]$
$[VP \rightarrow \bullet aux \ VP \ , \ 4]$
$[VP \rightarrow \bullet v \ NP \ , \ 4]$

S_5: $[VP \rightarrow v_6 \ \bullet NP \ , \ 4]$ 6 v

Earley Parsing - Example

The input: "$x =$ art adj n aux v **art** n \$"

$$
\begin{aligned}
S_5\text{:}\quad & [VP \rightarrow v_6 \ \bullet NP\ ,\ 4] \\
& [NP \rightarrow \bullet art\ adj\ n\ ,\ 5] \\
& [NP \rightarrow \bullet art\ n\ ,\ 5] \\
& [NP \rightarrow \bullet adj\ n\ ,\ 5]
\end{aligned}
$$

$$
\begin{aligned}
S_6\text{:}\quad & [NP \rightarrow art_7 \ \bullet adj\ n\ ,\ 5] \qquad\qquad 7 \quad art \\
& [NP \rightarrow art_7 \ \bullet n\ ,\ 5]
\end{aligned}
$$

Earley Parsing - Example

The input: "$x =$ art adj n aux v art **n** \$"

S_6: $[NP \rightarrow art_7 \ \bullet adj \ n \ , \ 5]$

$[NP \rightarrow art_7 \ \bullet n \ , \ 5]$

S_7: $[NP_9 \rightarrow art_7 \ n_8 \bullet \ , \ 5]$

8 n

9 $NP \rightarrow art_7 \ n_8$

Earley Parsing - Example

The input: "$x =$ art adj n aux v art n **$**"

S_7: $[NP_9 \rightarrow art_7\ n_8 \bullet ,\ 5]$

$[VP_{10} \rightarrow v_6\ NP_9 \bullet ,\ 4]$

$[VP_{11} \rightarrow aux_5\ VP_{10} \bullet ,\ 3]$

$[S_{12} \rightarrow NP_4\ VP_{11} \bullet ,\ 0]$

$[S' \rightarrow S\ \bullet\ \$,\ 0]$

10	$VP \rightarrow v_6\ NP_9$
11	$VP \rightarrow aux_5\ VP_{10}$
12	$S \rightarrow NP_4\ VP_{11}$

S_8: $[S' \rightarrow S\ \$ \bullet ,\ 0]$