# Concurrent Object Oriented Languages Synchronous Message Passing

https://wiki.cse.yorku.ca/course/6490A

## Communicating Sequential Processes (CSP)

C.A.R. Hoare. Communicating sequential processes. *Communications of the ACM*, 21(8):666-677, August 1978.



sir Charles Antony Richard (Tony) Hoare source: cs.ox.ac.uk

# Communicating Sequential Processes (CSP)

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CSP has static process creation.

```
[ name :: command \| \cdots \| name :: command ]
```

CSP uses synchronous message passing to communicate.

Receive command

name?pattern

Send command

name!expression

#### Communication in CSP

#### Question

What is the result of the following communication?

#### Answer

The variable x is assigned the value 2.

#### Communication in CSP

#### Question

What is the result of the following communication?

#### Answer

No communication takes place since the expression (1,2) does not match the pattern (3,x).

#### Conditional command

```
[ \mathtt{guard} 	o \mathtt{command} \ \square \ \cdots \ \square \ \mathtt{guard} \ 	o \mathtt{command} ]
```

#### guard

- Boolean expression
- receive command
- Boolean expression; receive command

#### Iteration command

```
*[ guard 	o command \square \cdots \square guard 	o command ]
```

#### guard

- Boolean expression
- receive command
- Boolean expression; receive command

Express a semaphore, named semaphore, and a process, named process, using that semaphore to protect its critical section in CSP.

Express the consumer-producer problem in CSP. The producer, named producer, produces the integers 1, ..., 100 and the consumer, named consumer, prints the integers it consumes. Both interact with the buffer, named buffer.

```
Let
```

```
reader(i) ::
    *[ scheduler!request();
        read();
        scheduler!done() ]

writer(i) ::
    *[ scheduler!request();
        write();
        scheduler!done() ]
```

Implement scheduler to solve the readers-writers problem.

#### What is wrong with

```
phil(i) ::
  *[ THINK;
     fork(i)!pickup(); fork((i+1) mod N)!pickup();
     EAT;
     fork(i)!putdown(); fork((i+1) mod N)!putdown()
fork(i) ::
  * [ phil(i)?pickup()
       → phil(i)?putdown()
     ☐ phil((i-1) mod N)?pickup()
       \rightarrow phil((i-1) mod N)?putdown() ]
```

The sieve of Eratosthenes is a simple, ancient algorithm for finding all prime numbers up to a specified integer.



Eratosthenes

source: world.mathigon.org

#### Processes:

- generator that generates 2, 3, ...
- sieve (i), for  $1 \le i \le n$ , where n is the number of primes to be generated (sieve (n) is defined differently).

```
sieve(0) ::
  n = 2:
  *[sieve(1)!n; n = n + 1]
sieve(i) ::
  sieve(i - 1)?p;
  print(p);
  *[ sieve(i - 1)?n
        \rightarrow [ n mod p == 0 \rightarrow skip
            \square n mod p != 0 \rightarrow sieve(i + 1)!n ]
sieve(100) ::
  sieve(99)?p; print(p)
```

# Assignment 1

Due: October 1

Presentations: October 8 and 13